

RYERSON UNIVERSITY
MTH 714 LAB#7 - SOLUTIONS

1. (a) Model for A_1 , A_2 but not A_3 : consider a set $U = \{a, b, c\}$ on which we have a binary relation $p = \{(a, a), (b, b), (c, c), (a, b), (b, a), (b, c), (c, b)\}$ (for example, one can visualize the model as a graph with three vertices a, b , and c , which has a loop at each vertex, and there are edges joining a and b and b and c , but no edge between a and c). Then, the binary relation is clearly reflexive and symmetric, but not transitive since

$$(a, b), (b, c) \in p \text{ but } (a, c) \notin p$$

- (b) Model for A_1 and A_3 , but not A_2 : take \mathbb{N} with the relation \leq . Clearly, $x \leq x$ is always true, and

$$\forall x \forall y \forall z (x \leq y \wedge y \leq z \rightarrow x \leq z)$$

but, e.g.

$$3 \leq 5 \text{ and } 5 \not\leq 3$$

which shows that A_2 fails.

- (c) Model for A_2 and A_3 but not A_1 : take $U = \{a, b, c\}$ with the binary relation $p = \{(a, b), (b, a), (b, c), (c, b), (a, c), (c, a)\}$. Again, we can visualize this structure as a graph with no loops which is a triangle with vertices a, b , and c . Then, the binary relation is symmetric and transitive but not reflexive.

2. (a) Yes. The formula translates as

$$F = \exists x \exists y \exists z (x < y \wedge z < y \wedge x < z \wedge z \not< x)$$

One example of such a triple of numbers is

$$x = 1, \quad y = 3, \quad z = 2$$

- (b) No. If such a triple of positive integers existed, it would mean that

$$y = x + 1, \quad y = z + 1, \quad z = x + 1, \quad x \neq z + 1.$$

However, the first two equalities yield $x = z$, but this is impossible according to the third equality.

- (c) Yes. The meaning of the formula in this structure is

$$F = \exists X \exists Y \exists Z (X \subseteq Y \wedge Z \subseteq Y \wedge X \subseteq Z \wedge Z \not\subseteq X)$$

and a triple of such subsets of \mathbb{N} is e.g.

$$X = \{1\}, \quad Y = \{1, 2\}, \quad Z = \{1, 2\}.$$

3. F Consider

$$F = \exists x \exists y \exists z [(p(x) \wedge q(x)) \wedge (\neg p(y) \wedge q(y)) \wedge \neg q(z)]$$

where p and q are two unary relation symbols. If

$$I \models F$$

we claim that $|I| \geq 3$. The reason for this is that, if x , y , and z are three elements that witness the truth of this formula, since

$$p(x) \text{ and } \neg p(y)$$

we must have $x \neq y$. Similarly, since

$$q(y) \text{ and } \neg q(z)$$

it must be the case that $y \neq z$. Also, we cannot have $x = z$, since $q(x)$ holds while $q(z)$ does not.

4. We will show that the formula is not valid by exhibiting an interpretation I in which the formula is false. Since the formula is in the form of an implication, we are looking for an interpretation I , which is a set with a binary relation on it, such that

$$\begin{aligned} I \models \forall x \forall y \forall z [p(x, x) \wedge (p(x, z) \rightarrow (p(x, y) \vee p(y, z)))] \\ I \models \neg \exists y \forall z p(y, z) \end{aligned}$$

The latter is equivalent to

$$I \models \forall y \exists z \neg p(y, z)$$

Consider the structure

$$I = (\mathbb{N}, \{\leq\})$$

Then,

$$\begin{aligned} I \models \forall x \forall y \forall z [x \leq x \wedge (x \leq z \rightarrow (x \leq y \vee y \leq z))] \\ I \models \neg \exists y \forall z (y \leq z) \end{aligned}$$

The former is true, since it is always the case that $x \leq x$, and if $x \leq z$ and $x \not\leq y$, then $y < x \leq z$ which yields $y \leq z$.

The latter formula is also true, since \mathbb{N} has the minimum element $y = 1$.